

# CS 603: Programming Languages

Lecture 27

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Department of Computer Science

University of Alabama

Joel Jones

# Overview

- Tracing append
- Introduction to Definite Clause Grammars

# Tracing Append

- See Handout

# Definite Clause Grammars

- Taken from:
  - <http://www.coli.uni-sb.de/~kris/prolog-course/html/node54.html>
- Prolog was originally intended for use in computational linguistics
- Definite Clause Grammars (DCGs) are syntactic sugar for implementing Context Free Grammars (CFGs) and other grammars

# Introduction to CFGs

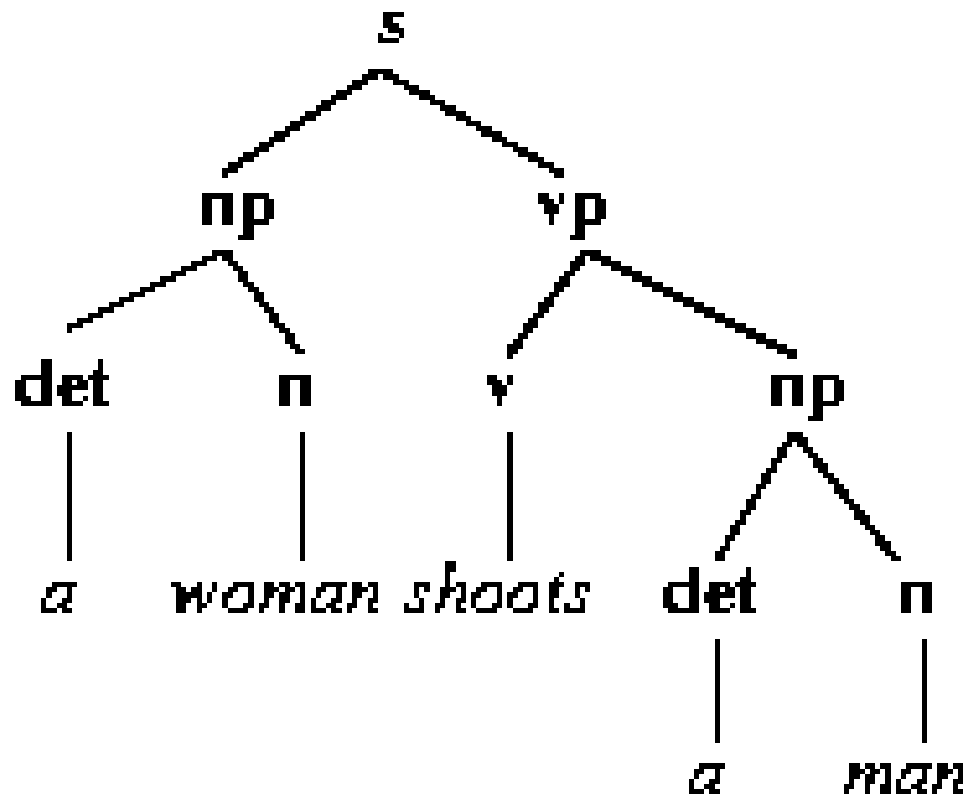
## Nonterminals

CFGs have only a single non-terminal on the left-hand side (LHS)

s	->	np vp
np	->	det n
vp	->	v np
vp	->	v
det	->	<i>a</i>
det	->	<i>the</i>
n	->	<i>woman</i>
n	->	<i>man</i>
v	->	<i>shoots</i>

Terminals

# Example Parse



# CFG Recognition Using append

```
s(Z) :- np(X), vp(Y), append(X, Y, Z).  
np(Z) :- det(X), n(Y), append(X, Y, Z).  
vp(Z) :- v(X), np(Y), append(X, Y, Z).  
vp(Z) :- v(Z).  
det([the]).  
det([a]).  
n([woman]).  
n([man]).  
v([shoots]).
```

This is nice—see all sentences  $s(X)$ .

But this is inefficient—trace  $s([a, woman, shoots, a, man])$ .

# CFG Recognition Using Difference Lists

$s(X, Z) :- \text{np}(X, Y), \text{vp}(Y, Z).$   
 $\text{np}(X, Z) :- \text{det}(X, Y), \text{n}(Y, Z).$   
 $\text{vp}(X, Z) :- \text{v}(X, Y), \text{np}(Y, Z).$   
 $\text{vp}(X, Z) :- \text{v}(X, Z).$   
 $\text{det}([\text{the} | W], W).$   
 $\text{det}([\text{a} | W], W).$   
 $\text{n}([\text{woman} | W], W).$   
 $\text{n}([\text{man} | W], W).$   
 $\text{v}([\text{shoots} | W], W).$

The **s** rule says: *I know that the pair of lists **x** and **z** represents a sentence if (1) I can consume **x** and leave behind a **y**, and the pair **x** and **y** represents a noun phrase, and (2) I can then go on to consume **y** leaving **z** behind, and the pair **y z** represents a verb phrase.*

This is efficient—trace  $s([\text{a}, \text{woman}, \text{shoots}, \text{a}, \text{man}], [])$ .



# Definite Clause Grammars

- DFGs allow us to write grammars without the ugliness of writing all those extra variables

```
s --> np, vp.  
np --> det, n.  
vp --> v, np.  
vp --> v.  
det --> [ the ].  
det --> [ a ].  
n --> [ woman ].  
n --> [ man ].  
v --> [ shoots ].
```

This is syntactic sugar for difference lists— `listing(s)`.

# Recursive Grammars

- Our original language only generated a finite number of sentences
- Add rules for conjunctions

$s \rightarrow np, vp.$

$np \rightarrow det, n.$

$vp \rightarrow v, np.$

$vp \rightarrow v.$

$det \rightarrow [the].$

$det \rightarrow [a].$

$n \rightarrow [woman].$

$n \rightarrow [man].$

$v \rightarrow [shoots].$

What happens if we add here?

$s \rightarrow s, conj, s.$

$conj \rightarrow [and].$

$conj \rightarrow [or].$

$conj \rightarrow [but].$

What happens if we add here and do  
 $s([woman, shoot])?$

We need to eliminate left recursion

# Left Recursion Eliminated

```
s --> simple_s.  
s --> simple_s, conj, s.  
simple_s --> np, vp.  
np --> det, n.  
vp --> v, np.  
vp --> v.  
det --> [the].  
det --> [a].  
n --> [woman].  
n --> [man].  
v --> [shoots].  
conj --> [and].  
conj --> [or].  
conj --> [but].
```