

CS603 Programming Language Organization

Lecture 23

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Department of Computer Science

Overview

- Questions
- Some final points about ML
- Introduction to Prolog

Welcome to the ML Realm

- Suppose we want to represent all inhabitants of a kingdom by their class:
 - King – simply himself
 - Peer – has a degree, territory and number
 - ex.: “the 7th Earl of Carlisle”
 - Knight or Peasant – has a name

How might we represent this in μ -Scheme?

The Realm in μ -Scheme

- "king"
- ("Earl" "Carlisle" 7) ("Duke" "Norfolk" 9)
- ("Knight" "Gawain") ("Knight" "Galahad")
- ("Peasant" "Jack Cade") ("Peasant" "Wat Tyler")
- But this doesn't work in ML!

Why?

More About ML

- Strings are supported with "abc", concatenation is caret, ^
 - example: "abc" ^ "def" yields "abcdef"
- Pattern matching can be done on datatypes
 - ```
datatype iTree = IntNode of int * iTree * iTree
 | Empty;
fun sum (IntNode(val,L,R)) = val + sum L + sum R
 | sum Empty = 0;
```

# Programs

- Write a datatype for representing a person

- ```
datatype person = King
                | Peer of ...
                | Knight of ...
                | Peasant of ... ;
```

Programs (cont.)

- Write a function `title` that behaves as follows:
`title King` yields "His Majesty the King"
`title Peer("Earl", "Carlisle", 7)` yields "The Earl of Carlisle"
`title (Knight "Gawain")` yields "Sir Gawain"
`title (Peasant "Joel")` yields "Joel"
- Write a function `sirs` that returns the name of all Knights in a list of persons, as a list of strings:
`sirs [King, Peasant("Joel"), Knight("Gawain")]`
yields ["Gawain"]

Solution

- ```
datatype person = King
 | Peer of string * string * int
 | Knight of string
 | Peasant of string;

fun title King = "His Majesty the King"
 | title (Peer(deg,terr,_)) = "The " ^ deg ^ " of " ^ terr
 | title (Knight name) = "Sir " ^ name
 | title (Peasant name) = name;

fun sirs [] = []
 | sirs ((Knight(s) :: ps)) = s :: (sirs ps)
 | sirs (_::ps) = sirs ps;
```



# ML Resources

- “Programming in Standard ML”
  - <http://www-2.cs.cmu.edu/~rwh/smlbook/offline.pdf>
- Source code from above book
  - <http://www-2.cs.cmu.edu/~rwh/smlbook/examples/>
- Moscow ML
  - <http://www.dina.dk/~setsoft/mosml.html>

# Prolog

- PROgramming in LOGic
- Algorithm = Logic + Control
- Logic = relation  $R(I,O)$  between input  $I$  and output  $O$
- Control = method of searching for  $O$  that satisfies  $R(I,O)$ , given input  $I$
- E.g. Find  $X$  and  $Y$  such that
  - $3*X+2*Y=1$
  - $X-Y=4$
- E.g. find array  $B$  such that
  - elements in  $B$  are the same as those in  $A$
  - elements of  $B$  are in non-descending order

# What is Prolog

- Prolog is a ‘typeless’ language with a very simple syntax.
- Prolog is declarative: you describe the relationship between input and output, not how to construct the output from the input (“specify what you want, not how to compute it”)
- Prolog uses a subset of first-order logic

# Classical First-Order Logic

- simplest form of logical statements is an atomic formula. e.g.
  - `man(tom)`
  - `woman(mary)`
  - `married(tom,mary)`
- More complex formulas can be built up using logical connectives:  $\wedge, \vee, \neg, \rightarrow, \forall X, \exists X$

Pair Up:

- Define each of these symbols

# Examples of First Order Logic

- $\text{smart}(\text{tom}) \vee \text{dumb}(\text{tom})$
- $\text{smart}(\text{tom}) \vee \text{tall}(\text{tom})$
- $\neg \text{dumb}(\text{tom})$
- $\exists X \text{ married}(\text{tom}, X)$
- $\exists X \text{ loves}(\text{tom}, X)$
- $\exists X [\text{married}(\text{tom}, X) \wedge \text{female}(X) \wedge \text{human}(X)]$
- $\text{rich}(\text{tom}) \vee \neg \text{smart}(\text{tom})$
- $\exists X \text{ mother}(\text{john}, X)$
- $\forall X \forall Y [\text{mother}(\text{john}, X) \wedge \text{mother}(\text{john}, Y) \rightarrow Y=X]$

**Note:**  $A \rightarrow B \equiv B \vee \neg A$

# Horn Rules

- Logic programming is based on formulas called Horn rules. These have the form

$$\forall x_1, \dots, x_k [A \leftarrow B_1 \wedge B_2 \wedge \dots \wedge B_j]$$

- Examples:

$$\forall X, Y [A(X) \leftarrow B(X, Y) \wedge C(Y)]$$

$$\forall X [A(X) \leftarrow B(X)]$$

$$\forall X [A(X, d) \leftarrow B(X, e)]$$

$$A(c, d) \leftarrow B(d, e)$$

$$\forall X A(X)$$

$$\forall X A(X, d)$$

$$A(c, d)$$

# Horn Rules (cont.)

- Note that atomic formulas are also Horn rules, often called facts.
- A set of Horn rules is called a Logic Program.

# Logical Inference with Horn Rules

- Logic programming is based on a simple idea: From rules and facts, derive more facts.
- Example 1. Given the facts and rules:
  1. A
  2. B
  3. C
  4.  $E \leftarrow A \wedge B$
  5.  $F \leftarrow C \wedge E$
  6.  $G \leftarrow E \wedge F$
- From 1, derive E; from 2, derive F; from 3, derive G.



# Logical Inference

- Example 2: Given these facts:
  - `man(plato)`
  - `man(socrates)`
- and this rule:
  - $\forall X [\text{mortal}(X) \leftarrow \text{man}(X)]$
- derive:
  - `mortal(plato), mortal(socrates).`