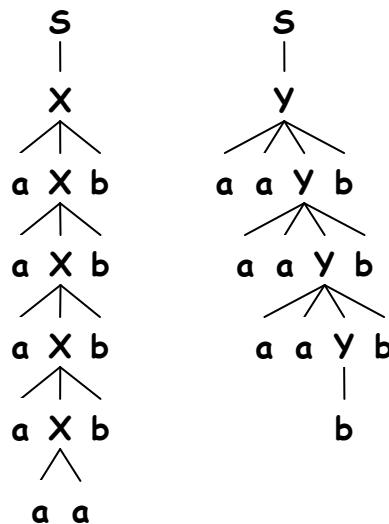
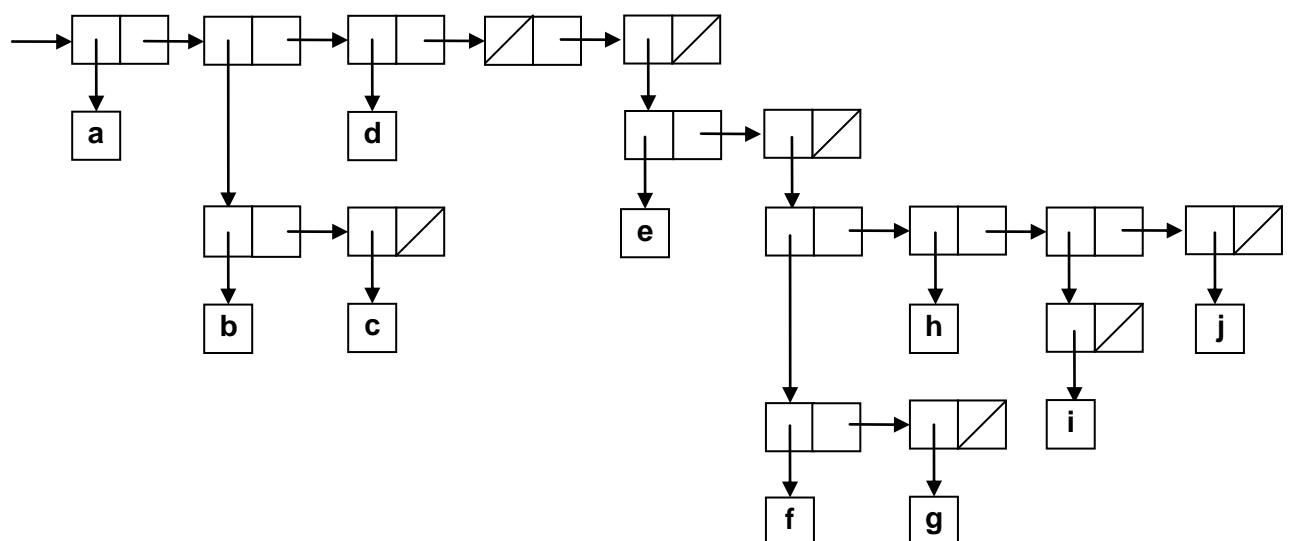


1. Show that this grammar is ambiguous by drawing two parse trees that generate the same string.

$$\begin{aligned} S &\rightarrow X \mid Y \\ X &\rightarrow aXb \mid aa \\ Y &\rightarrow aaYb \mid b \end{aligned}$$


2. Draw a diagram that illustrates the internal representation of this S-expression:
 $(a (b c) d () (e ((f g) h (i) j)))$



3. Write an unambiguous context-free grammar that generates S-expressions that contain only the symbol x. For example, $(x (x x) x () (x ((x x) x (x) x)))$. Your production rules may use concatenation and |, but not other extended BNF operations.

$$\begin{aligned} S &\rightarrow (L) \mid x \\ L &\rightarrow SL \mid \epsilon \end{aligned}$$

4. Suppose we give Impcore a new primitive function `&&` which implements binary short-circuit AND. So the expression $(\&& x y)$ in Impcore should be equivalent to $x \& y$ in C or C++ or Java. Define the `&&` function by completing these two Impcore-style natural operational semantics rules.

$\phi(f) = \text{PRIMITIVE}(\&\&)$ $\langle e_1, \xi, \phi, \rho \rangle \Downarrow \langle v_1, \xi', \phi, \rho' \rangle$ $v_1 \neq 0$ $\langle e_2, \xi', \phi, \rho' \rangle \Downarrow \langle v_2, \xi'', \phi, \rho'' \rangle$ <hr/> $\langle \text{APPLY}(f, e_1, e_2), \xi, \phi, \rho \rangle \Downarrow \langle v_2, \xi'', \phi, \rho'' \rangle$	$\phi(f) = \text{PRIMITIVE}(\&\&)$ $\langle e_1, \xi, \phi, \rho \rangle \Downarrow \langle v_1, \xi', \phi, \rho' \rangle$ $v_1 = 0$ <hr/> $\langle \text{APPLY}(f, e_1, e_2), \xi, \phi, \rho \rangle \Downarrow \langle 0, \xi', \phi, \rho' \rangle$
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5. Suppose we give μ Scheme a new primitive function `||` which implements binary short-circuit OR. So the expression $(|| x y)$ in μ Scheme should be equivalent to $x || y$ in C or C++ or Java. Define the `||` function by completing these two μ Scheme-style natural operational semantics rules.

$\langle e, \rho, \sigma \rangle \Downarrow \langle \text{PRIMITIVE}(), \sigma_1 \rangle$ $\langle e_1, \rho, \sigma_1 \rangle \Downarrow \langle v_1, \sigma_2 \rangle$ $v_1 \neq \text{BOOL}(\#f)$ <hr/> $\langle \text{APPLY}(e, e_1, e_2), \rho, \sigma \rangle \Downarrow \langle v_1, \sigma_2 \rangle$	$\langle e, \rho, \sigma \rangle \Downarrow \langle \text{PRIMITIVE}(), \sigma_1 \rangle$ $\langle e_1, \rho, \sigma_1 \rangle \Downarrow \langle v_1, \sigma_2 \rangle$ $v_1 = \text{BOOL}(\#f)$ <hr/> $\langle e_2, \rho, \sigma_2 \rangle \Downarrow \langle v_2, \sigma_3 \rangle$ <hr/> $\langle \text{APPLY}(e, e_1, e_2), \rho, \sigma \rangle \Downarrow \langle v_2, \sigma_3 \rangle$
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6. Write a μ Scheme function (`diagonal M`) where M is a square matrix stored as a list of row lists. It should return a list of the main diagonal elements of M . Example: `(diagonal '((a b c d) (e f g h) (i j k l) (m n o p)))` returns `(a f k p)`.

$$\begin{bmatrix} a & b & c & d \\ e & f & g & h \\ i & j & k & l \\ m & n & o & p \end{bmatrix}$$

```
(define diagonal (M) (if (null? M) '()
  (cons (car (car M)) (diagonal (map cdr (cdr M))))))
```

7. Write a μ Scheme function (`scan op id L`) where op is a binary function, id is the identity value, and L is a list. It should return a list of the values obtained by folding op across each possible prefix of L . Example: `(scan + 0 '(2 3 5 7 11))` returns `(0 2 2+3 2+3+5 2+3+5+7 2+3+5+7+11) = (0 2 5 10 17 28)`. [You may assume that op is an associative operation.]

```
(define scan (op id L) (if (null? L) (cons id '())
  (cons id (map ((curry op) (car L)) (scan op id (cdr L))))))
```

8. The Impcore function below shows an inefficient way to compute Fibonacci numbers. Note that $(\text{fib } 0) = (\text{fib } 1) = 1$, $(\text{fib } 2) = 2$, $(\text{fib } 3) = 3$, $(\text{fib } 4) = 5$, $(\text{fib } 5) = 8$, etc. Write an equivalent function so that $(\text{fib } n)$ runs in $O(n)$ time. Hint: use a helper function that "remembers" the previous two values.

```
(define fib (n) (if (<= n 1) 1
  (+ (fib (- n 2)) (fib (- n 1)))))
```

```
(define fib (n) (helper n 1 1))
```

```
(define helper (n x y) (if (<= n 1) y
  (helper (- n 1) y (+ x y))))
```

9. Complete the μ Scheme function Stack so that the print statements in the client code below will produce the given output.

(define Stack () (let ((L '())) (lambda (m) (if (= m 'isEmpty) (null? L) (if (= m 'top) (car L) (if (= m 'pop) (set L (cdr L)) (if (= m 'push) (lambda (v) (set L (cons v L))) 'error)))))))	(val A (Stack)) (val B (Stack)) (val k 0) (while (< k 10) (begin ((A 'push) k) ((B 'push) (+ k 1)) (set k (+ k 2))))) (while (not (A 'isEmpty)) (begin (print (A 'top)) (A 'pop))) (while (not (B 'isEmpty)) (begin (print (B 'top)) (B 'pop))))	8 6 4 2 0 9 7 5 3 1
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10. Using either C or C++ or Java, write a definition for a Stack abstract data type represented as a linked list of ints. Provide these operations with the same functionality as in the preceding problem: isEmpty, top, pop, push.

```
// Java
class Node {
    int value;
    Node next;
    Node(int v, Node n) { value=v; next=n; }
}
class Stack {
    Node head;
    Stack() { head = null; }
    boolean isEmpty() { return head==null; }
    int top() { return head.value; }
    void pop() { head = head.next; }
    void push(int v) { head = new Node(v, head); }
}
```